

Accelerating Dynamic Programming algorithms using Massively Parallel Processors

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Why?

Massively Parallel Processing now **matters**.

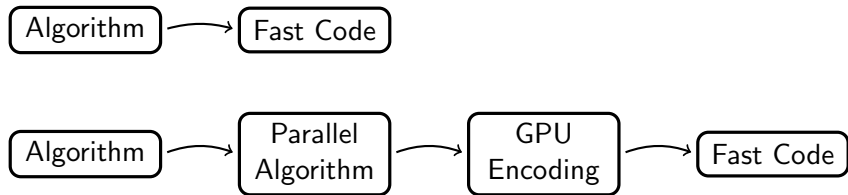
GPUs now power **3 of the top 5** supercomputers

Installed base ranges from supercomputers to netbooks

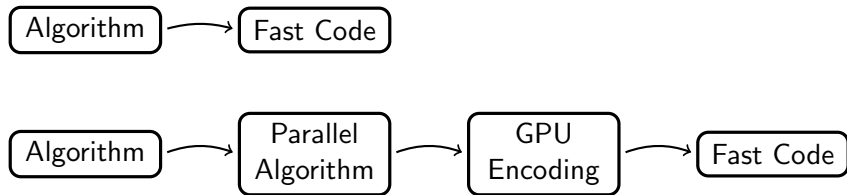
Mapping problems is hard



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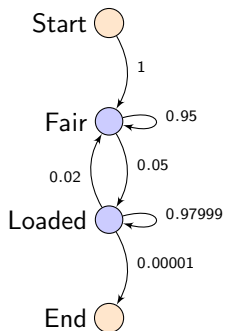


Mapping problems is hard



General-purpose automated parallelisation is difficult.
Instead, lets focus on **domains**.

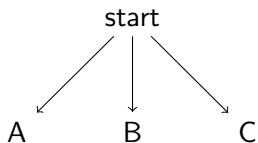
Progress so far



HMMingbird

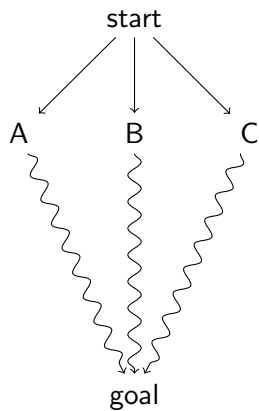
- A Hidden Markov Model compiler
- Targets GPUs
- Using a fixed set of algorithms

Optimisation Problems



Choices

Optimisation Problems



Choices

Optimal Substructure

Edit distance

Minimum operations required to transform one string to another.

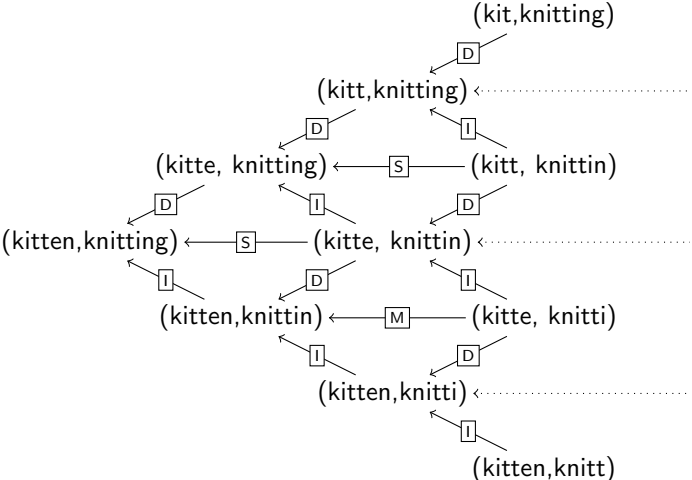
e.g

kitten \implies kitteng (I)

kitteng \implies kitting (S)

kitting \implies knitting (I)

Graph problem



Recursion

$$f(x, y) = \begin{cases} x & \text{if } y = 0 \\ y & \text{if } x = 0 \\ f(x - 1, y - 1) & \text{if } s[x] = t[y] \\ \min(f(x - 1, y), \\ f(x, y - 1), \\ f(x - 1, y - 1)) + 1 & \text{otherwise} \end{cases}$$

Partitions

| | | k | n | i | t | t | i | n | g |
|----------|---|----------|----------|----------|----------|----------|----------|----------|----------|
| | 0 | ← 1 | ← 2 | ← 3 | ← 4 | ← 5 | ← 6 | ← 7 | ← 8 |
| k | 1 | ← 0 | ← 1 | ← 2 | ← 3 | ← 4 | ← 5 | ← 6 | ← 7 |
| i | 2 | ← 1 | ← 1 | ← 1 | ← 2 | ← 3 | ← 4 | ← 5 | ← 6 |
| t | 3 | ← 2 | ← 2 | ← 2 | ← 1 | ← 2 | ← 3 | ← 4 | ← 5 |
| t | 4 | ← 3 | ← 3 | ← 3 | ← 2 | ← 1 | ← 2 | ← 3 | ← 4 |
| e | 5 | ← 4 | ← 4 | ← 4 | ← 3 | ← 3 | ← 2 | ← 3 | ← 4 |
| n | 6 | ← 5 | ← 4 | ← 5 | ← 4 | ← 4 | ← 4 | ← 2 | ← 3 |

Partitions

| | | k | n | i | t | t | i | n | g |
|----------|---|----------|----------|----------|----------|----------|----------|----------|----------|
| | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| k | 1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| i | 2 | 1 | 1 | 1 | 2 | 3 | 4 | 5 | 6 |
| t | 3 | 2 | 2 | 2 | 1 | 2 | 3 | 4 | 5 |
| t | 4 | 3 | 3 | 3 | 2 | 1 | 2 | 3 | 4 |
| e | 5 | 4 | 4 | 4 | 3 | 3 | 2 | 3 | 4 |
| n | 6 | 5 | 4 | 5 | 4 | 4 | 4 | 2 | 3 |

Partition function

Describe each partition based on the parameters e.g

$$P_f(x, y) = x + y$$

Proving suitability of a partition function

Consider each recursive call in turn using descent functions.

$$f(d_x(x, y), d_y(x, y))$$

e.g

$$f(x - 1, y - 1)$$

Verify the recursive equation satisfies the partition function:

$$\forall_{x,y} P_f(x, y) > P_f(d_x(x, y), d_y(x, y))$$

e.g

$$\forall_{x,y} x + y > (x - 1) + (y - 1)$$

Termination

Partition base case is when $P_f(x, y) = 0$:

The base case of the recursion needs to satisfy this.

We must also check the recursion never jumps out of bounds.

Further work

Tie work into higher-level domains for Bioinformatics
(e.g Hidden Markov Models)
Permit more sophisticated recursions